

# Data-driven Modeling and Coordination of Large Process Structures\*

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**Abstract.** In the engineering domain, the development of complex products (e.g., cars) necessitates the coordination of thousands of (sub-)processes. One of the biggest challenges for process management systems is to support the modeling, monitoring and maintenance of the many interdependencies between these sub-processes. The resulting process structures are large and can be characterized by a strong relationship with the assembly of the product; i.e., the sub-processes to be coordinated can be related to the different product components. So far, sub-process coordination has been mainly accomplished manually, resulting in high efforts and inconsistencies. IT support is required to utilize the information about the product and its structure for deriving, coordinating and maintaining such *data-driven process structures*. In this paper, we introduce the COREPRO framework for the data-driven modeling of large process structures. The approach reduces modeling efforts significantly and provides mechanisms for maintaining data-driven process structures.

## 1 Introduction

Enterprises are increasingly demanding IT support for their business processes. One challenge emerging in this context is to coordinate the execution of large and long-running processes (e.g., related to car development). Engineering processes, for instance, often consist of numerous concurrently executed, interdependent sub-processes. The reasons for this fragmentation are manifold: Typically, these sub-processes are related to different (data) objects (e.g., product components), enacted by different organizational units (e.g., dealing with the *testing* or *releasing* of single components), and controlled by different IT systems. We denote such correlated sub-processes as *process structure*.

These process structures have in common that changes (e.g., removing a sub-process or adding a dependency between sub-processes) as well as real-world exceptions (e.g., abnormal termination of a sub-process) occur frequently and may

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affect not only single sub-processes but also the whole process structure [1, 2]. Consequently, IT support must flexibly cover modeling, enactment and maintenance of process structures and assure their consistency. Even the modeling of process structures constitutes a challenging task since these structures usually comprise hundreds up to thousands of sub-processes (and sub-process dependencies). Doing this manually often results in errors or inconsistencies leading to bad process performance and high process costs.

To cope with these challenges, we have to better understand the dependencies between sub-processes. Case studies we conducted in the automotive industry [1, 3] have revealed that the dependencies between the different sub-processes of a process structure typically base on the assembly of the product to be manufactured. As an example for a *product structure* (or configuration structure) consider the total electrical system in a modern car which consists of up to 300 interconnected components. To verify the functionality of the total system, several sub-processes (e.g., *testing* and *release*) have to be executed for each electrical (sub-)component. Interestingly, the technical relations between the different product components indicate sub-process dependencies; i.e., the relation between two components leads to dependencies between sub-processes modifying these components. We use the notion of *data-driven process structures* to characterize process structures, which are prescribed by respective data structures. Fig. 1 presents an example for a data-driven process structure. The strong relationship between data structures and process structures (e.g., the relations between the S-Class and the Navigation object leads to respective sub-process dependencies) also implies that a changed data structure (e.g., total electrical system for another car series without navigation) leads to a different process structure.

Our goal is to reduce modeling efforts for data-driven process structures by increasing model reusability and maintainability. In the automotive domain, for instance, we can benefit from the upcoming standardization of develop-

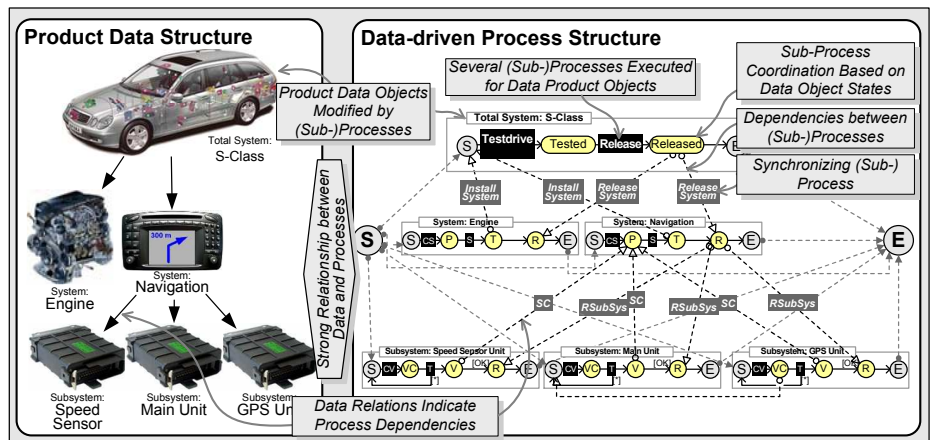


Fig. 1. Example for a Data Structure and a Related Data-driven Process Structure

ment processes driven by quality frameworks like CMMI (Capability Maturity Model Integration) or engineering guidelines (e.g., [4]). This leads to standardized processing of objects (e.g., the testing process for the speed sensor is independent from the car series it is built in), which can be utilized to increase reuse of process models and to reduce modeling efforts. In order to benefit from the standardization, a loose coupling of data structures and process structures is required. In particular, three issues arise:

1. How to describe the processing of single objects, i.e., the relationship between an object and its modifying sub-processes?
2. How to describe the processing of the overall data structure, i.e., the dependencies between the sub-processes in relation to the different data objects?
3. How to automatically derive a proper process structure that can be *customized* by the underlying data structure, i.e., different data structures lead to different process structures?

So far, there exists no IT support for configuring a process structure based on a given data structure. IT systems used in industry, such as product data management systems or workflow management tools currently lack an integration of data and processes [1]. Approaches from academia, such as data-centered process paradigms [5–7] also do not fully address the aforementioned issues. Instead, users have to manually define the requested process structure for each given data structure. This often leads to inflexible process structures and generates high efforts for coordinating and maintaining them. In this paper, we introduce the modeling component of the COREPRO approach, which aims at an intuitive and product-related integration of data and (sub-)processes. In particular, COREPRO enables

- the data-driven specification of process structures at the model level
- the automated creation of process structures based on given data structures
- the data-driven adaptation of process structures to deal with real-world changes.

We utilize the life cycles of objects (i.e., the sequence of states an object goes through during its lifetime) for enabling data-driven modeling and coordination of process structures. State transitions within a life cycle take place when a sub-process is enacted for the related object [8–11]. According to the relations between objects, we connect the life cycles of these objects. The concept thereby enables the automated derivation of the process structure for a given data structure.

The remainder of this paper is structured as follows: Section 2 characterizes the relationship between data and process structures, and introduces our approach for describing them. Section 3 shows how to model and instantiate (product) data structures. Section 4 deals with the data-driven creation and change of large process structures. Section 5 illustrates the practical benefits of the COREPRO approach. Section 6 discusses related work, and Section 7 concludes with a summary and outlook.

## 2 Overview of the Approach

IT support for the modeling and change of data-driven process structures must meet four major requirements. First, it must enable the definition of the (product) data structure, i.e., its objects and their relations. Second, with each (data) object a set of sub-processes for processing this object and for transforming its state has to be associated. Third, sub-process dependencies have to be defined based on the object relations. For example, a sub-process for the **Navigation** object shall be not started before having finished the sub-processes of the related subsystems (cf. Fig. 1). Fourth, the concepts must enable the automated creation of a data-driven process structure.

The COREPRO modeling framework meets these requirements. In order to enable reuse and to reduce modeling efforts, COREPRO distinguishes between the model and the instance level when creating data-driven process structures (cf. Fig. 2). We allow defining a domain specific *data model* consisting of object and relation types (cf. Step 1a in Fig. 2). Such a data model can then be instantiated to create specific *data structures* (e.g., a *bill of material*) consisting of objects and relations (cf. Step 1b in Fig. 2). While the definition of a data model requires profound domain knowledge, the instantiation can be done by users.

Further, process experts describe the dynamic aspects of each object type by modeling *object life cycles* (OLC). An OLC defines the coordination of sub-processes associated with a particular object type (cf. Step 2a in Fig. 2). A sub-process is an autonomous process (or activity). In COREPRO an OLC is mapped

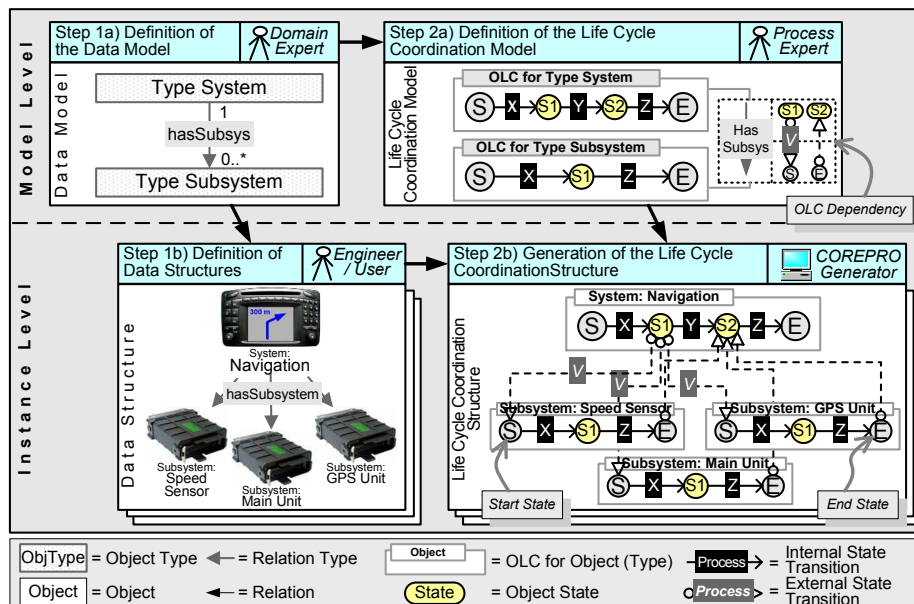


Fig. 2. Overall Concept of the COREPRO Modeling Approach

to a state transition system whose states correspond to object states and whose (internal) state transitions are triggered when associated sub-processes (which are modifying the object) are completed. In Fig. 2, the OLC for object type **System** (Step 2a), for example, goes through state **S1** followed by state **S2**. The *internal state transition* from **S1** to **S2** takes place when finishing sub-process **Y**. Altogether, an OLC constitutes an integrated view on a particular object and on the sub-processes manipulating this object.

Defining the dynamic aspects of each object type by modeling its OLC is only half of the story. We also have to deal with the many dependencies existing between the sub-processes associated with different objects types. Such dependencies might describe, for example, that every system (**Engine** and **Navigation**) must have reached a certain state before the **Testdrive** sub-process for the **S-Class** can be started (cf. Fig. 1). Consequently, a sub-process dependency can be seen as a synchronization link between the states of concurrently enacted OLCs.

In COREPRO, we specify such sub-process dependencies by defining *external state transitions*, which connect states of different OLCs. Like an internal state transition within an OLC, an external state transition can lead to the enactment of a sub-process. In Fig. 2, for example, the external state transition between the OLCs of **Type System** and **Type Subsystem** is associated with sub-process **V** (cf. Step 2a). Further, external state transitions are mapped to relation types.

In COREPRO, the OLCs for every object type and the external state transitions for every relation type form the *Life Cycle Coordination Model* (LCM) (cf. Fig. 2, Step 2a). Consequently, the LCM describes the dynamic aspects of the whole data model and constitutes the blueprint for creating the data-driven process structure.

On instance level, the *life cycle coordination structure* (LCS) describes the process structure for a particular data structure. While data model, data structure, and LCM are created manually, the LCS can be automatically generated based on these ingredients (cf. Step 2b in Fig. 2). The LCS includes an OLC for every object in the data structure. Likewise, for each relation in the data structure, external state transitions are inserted to the LCS. For example, for every **hasSubsystem** relation in the data structure from Fig. 2 (Step 1b), the associated external state transitions (with the associated sub-process **V**) are inserted. The result is an enactable process structure describing the dynamic aspects of the given data structure. Further details are presented in the following sections.

### 3 Modeling and Instantiation of Data Structures

In COREPRO, a domain specific data structure establishes the basis for creating data-driven process structures. COREPRO enables the definition of dynamic aspects of objects and relations at model level. Therefore, the definition of a data model may consist of object and relation types (cf. Fig. 3a). Based on this, data structures can then be created by instantiating specific objects and relations (cf. Fig. 3b).

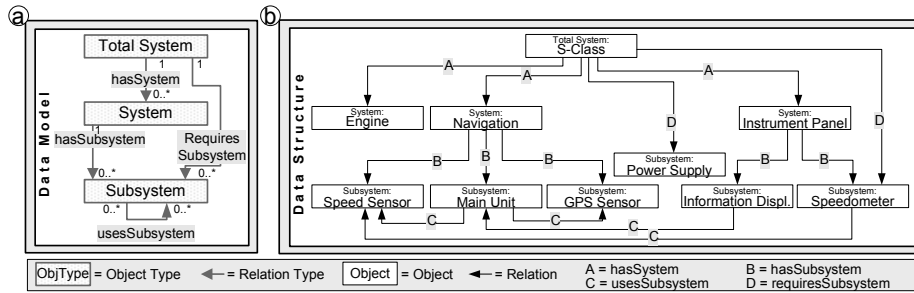


Fig. 3. a) Data Model and b) a Possible Instantiation (Data Structure)

### 3.1 Creation of a Data Model

Generally, a *data meta model* provides the constructs for describing a *data model* [12–14]. In COREPRO, we use a simple data meta model, comprising *object* and *relation types*<sup>1</sup>. An object type represents, for example, an abstract or physical product component that is part of the logical structure of a car. A relation type expresses a single relationship between two object types (including cardinality). A data model comprises object and relation types, and describes *how* objects are structured in a specific domain (cf. Fig. 3a). Based on the restrictions of the data model, several *data structures* (i.e., instances of the data model) can be created, such as product structures for different car series.

Generally, multiple relation types can be defined between two object types. Further, we allow defining recursive relation types, which can be used to realize relations between objects of the same object type on instance level.

**Definition 1 (Data Model).** Let  $\mathcal{T}$  be the set of all object types and let  $\mathcal{R}$  be the set of all relation types. Then: A data model is a tuple  $dm = (OT, RT)$  where

- $OT \subseteq \mathcal{T}$  comprises a set of object types
- $RT \subseteq OT \times \mathcal{R} \times OT$  comprises a set of binary relation types defined for object types
- $card : RT \mapsto \mathbb{N}_0 \times \mathbb{N}_0$  with  $card(ot_1, rt, ot_2) = (min_{rt}, max_{rt})$  assigns to each relation type  $rt \in RT$  a minimal and maximal cardinality.

### 3.2 Creation of the Data Structure

A data structure contains *objects* and *relations*, which constitute instances of the object and relation types defined by the data model. In Fig. 3b, for example, the **Total System** type is instantiated once, while the **System** type is instantiated three times (**Engine**, **Navigation**, and **Instrument Panel**). The cardinalities associated with relation types restrict the number of concrete relations between

<sup>1</sup> The data meta model neglects descriptive object attributes since they do not influence the generation of the requested process structure.

objects. Accordingly, a **Subsystem** object can only be related to one **System** object via the **hasSubsystem** relation (cf. Fig. 3b). Note that the data structure from Fig. 3b is a simplified example. In Sect. 5 we indicate, that such a data structure may comprise a high number of instantiated objects in practice.

**Definition 2 (Data Structure).** *Let  $dm = (OT, RT)$  be a data model. Then: A data structure created from  $dm$  is a tuple  $ds = (O, R)$  with*

- $O$  is a set of objects where each object  $o \in O$  is associated with an object type  $objtype(o) \in OT$
- $R \subseteq O \times RT \times O$  is a set of object relations meeting the cardinality constraints defined by the data models. Each relation  $r = (o_1, rt, o_2) \in R$  has a relation type  $rt = reltype(r)$  with  $(objtype(o_1), rt, objtype(o_2)) \in RT$ .

## 4 Integration of Data and Processes

After having defined how a data structure is modeled, we need to specify its relationship to the process structure. To allow for reuse, we describe this relationship at the model level; i.e., we define the dynamic aspects for the data model and translate them to the instance level afterwards. Therefore, an *object life cycle* (OLC) describes the dynamic aspects of an object type and an OLC dependency defines the dynamic aspects of a relation type. The *life cycle coordination model* (LCM) comprises the OLCs for every object type and the OLC dependencies for every relation type. Consequently, the LCM describes the dynamic aspects of the whole data model. On instance level, in turn, the LCM constitutes the basis for creating *life cycle coordination structures* (LCS) for given data structures. Altogether, the LCS defines the dynamic aspects of the underlying data structure and represents the data-driven process structure.

### 4.1 Modeling of the Dynamic Aspects of Single Object Types

Every object type is associated with an OLC, which constitutes a labeled transition system describing *object states* and *internal state transitions* (cf. Fig. 4). An internal state transition can be associated with a sub-process modifying the object (and thus inducing a state change). This sub-process becomes enacted when the source state of the transition is enabled. After having executed it, the source state of the transition is disabled and the target state becomes enabled. Hence, the definition of the OLC constitutes the mediator for associating stateful objects with modifying sub-processes<sup>2</sup>.

To realize non-deterministic processing, the definition of conditional (i.e., non-deterministic) internal state transitions is possible in COREPRO. All internal state transitions with same OLC state as source are associated with the same

<sup>2</sup> Stateless objects are also supported. Their OLCs include one internal state transition (with an associated sub-process) from the start to the end state.

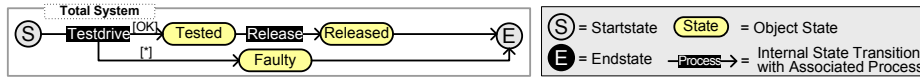


Fig. 4. Object Life Cycle with Conditional State Transitions

sub-process, but can be bound to (different) sub-process results, i.e., exit codes (e.g., *finished with errors*). Depending on the concrete sub-process result, always one internal state transition is triggered at runtime. In Fig. 4, for example, the result of the *Testdrive* sub-process determines, whether the state *Tested* or *Faulty* is enabled. The default transition to state *Faulty* (indicated by \*) becomes activated in case of uncovered sub-process results. A conditional internal transition may also be used for modeling loops within an OLC.

**Definition 3 (Object Life Cycle).** Let  $dm = (OT, RT)$  be a data model and let  $ot \in OT$  be an object type. Then: The object life cycle of  $ot$  is a tuple  $olc = (P, V, TS)$  where

- $P$  is a set of sub-processes that can be applied to instances of object type  $ot$  and  $V$  is a set of possible sub-process results ( $\sigma : P \mapsto \mathcal{P}(V)$  with  $\sigma(p) \subseteq V$  is the set of possible results defined for sub-process  $p$ )
- $TS = (S, T, s_{start}, s_{end})$  is a labeled transition system, where
  - $S$  is the set of states that can be reached by objects of type  $ot$
  - $T \subseteq S \times (P \times V) \times S$  is a set of *internal state transitions* with
    - \*  $t = (s, (p, v), s') \in T, \Rightarrow v \in \sigma(p)$ ; i.e., a state transition  $t$  is triggered by the completion of a sub-process  $p$  with particular result  $v$
    - \*  $\forall t_i = (s_i, (p_i, v_i), s'_i) \in T, i = 1, 2$  with  $t_1 \neq t_2$  and  $s_1 = s_2, \Rightarrow p_1 = p_2 \wedge v_1 \neq v_2$ ; i.e., if there are several state transitions with same source state  $s$ , all of them will be associated with the same sub-process  $p$ . The concrete target state is determined based on the sub-process result. In case of non-deterministic state transitions, there is a default transition that will be chosen if the associated sub-process delivers a result not covered by the other transitions.
  - $s_{start} \in S$  is the initial state and  $s_{end} \in S$  is the final state of the transition system;  $s_{start}$  is the only state without incoming transitions and  $s_{end}$  is the only state without outgoing transitions.

Let  $OLC$  be the set of all object life cycles. For  $olc \in OLC$ ,  $s_{start}(olc)$  denotes the start and  $s_{end}(olc)$  the end state of the respective transition system.

#### 4.2 Modeling of the Dynamic Aspects of the Data Model

Modeling the dynamic aspects of single object types is only one part of the challenge. To define the processing of the whole data model, we also have to specify the dynamic aspects of relation types. Relation Types are associated with *OLC dependencies* which synchronize the OLCs of related objects. An *OLC dependency* comprises several *external state transitions* between the states of the



dependent OLCs. The resulting structure is denoted as *life cycle coordination model* (LCM). A LCM describes the dynamic aspects of the data model by integrating the OLCs associated with object types as well as the OLC dependencies associated with relation types (Fig. 5a presents the LCM for the data model from Fig. 3a).

Like internal state transitions (within an OLC), an external state transition can be associated with a sub-process. As an example, consider the OLC dependency of the relation type `hasSystem` in Fig. 5a. This dependency consists of two external state transitions, which synchronize (1) the `start` state of the `Total System` OLC with the `Tested` state of the `System` OLC and (2) the `Release` state of the `System` OLC with the `Release` state of the `Total System` OLC. The sub-process associated with the external state transition (e.g., sub-process `InstallComponent`) can be considered as *synchronizing (sub-)process*, which operates on both related object types.

**Definition 4 (OLC Dependency).** Let  $olc_i = (P_i, V_i, TS_i), i = 1, 2$  be two different object life cycles with  $TS_i = (S_i, T_i, s_{start}, s_{end})$  (cf. Def. 3). Then: An OLC dependency between  $olc_1$  and  $olc_2$  is a tuple  $olc_{Dep} = (Id, P, EST)$  where

- $Id$  is the identifier of the dependency
- $P$  is a set of sub-processes that can operate on both object types
- $EST$  is a set of *external state transitions* with
 
$$est = (s, p, s') \in EST \Leftrightarrow (s \in S_1 \wedge s' \in S_2) \vee (s' \in S_1 \wedge s \in S_2).$$

$OLC_{DEP}$  denotes the set of all OLC dependencies. For an OLC dependency  $olc_{Dep} \in OLC_{DEP}$ , let  $est(olc_{Dep})$  denote the set of related external state transitions.

It is important to mention that internal and external state transitions differ in their operational semantics (cf. Table 1). During runtime, the concurrent processing of different objects is required to enhance process efficiency (e.g., by supporting concurrent engineering techniques). In COREPRO, this is realized by concurrently enacting different OLCs while avoiding concurrency within an OLC<sup>3</sup>. Both, internal and external state transition become activated (i.e., their sub-processes are started) when the source state of the transition is entered. While the completion of the sub-process of an internal state transition induces

<sup>3</sup> Concurrently activated states are not allowed within one OLC since this has not been a requirement in our case studies.

**Table 1.** Classification of State Transitions in COREPRO

Type	Meaning	Operational Semantics
Internal State Transition	Connects two states within one OLC	Fires after sub-process execution dependent on sub-process result
External State Transition	Connects two states from different OLCs	Fires after sub-process execution
Direct State Transition	Connects LCS start state with start state of an OLC (end states accordingly)	Fires immediately

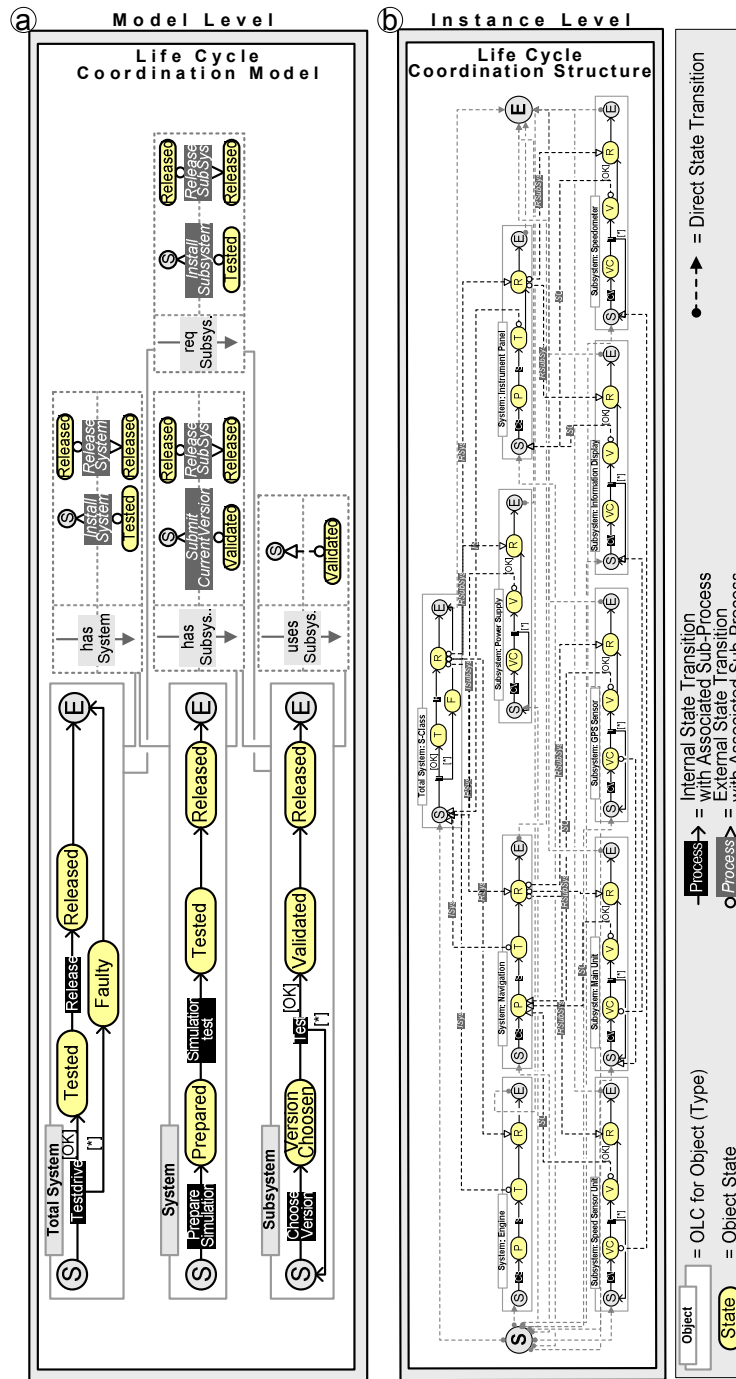


Fig. 5. Example for a) LCM and b) generated LCS

the deactivation of the source state and the activation of the target state, the completion of the sub-process of an external state transition does not imply any state change in the source OLC.

Further, the target state is activated if and only if the sub-processes of (1) one incoming internal state transition and (2) all incoming external transitions are fired. This rule allows for concurrently activated states within different OLCs of an LCS, while it prevents concurrently activated states within a single OLC. Due to the lack of space, we omit a formal specification of the operational semantics of internal and external state transitions.

**Definition 5 (Life Cycle Coordination Model).** *Let  $dm = (OT, RT)$  be a data model and let  $P$  be a set of sub-processes. Then: The life cycle coordination model associated with  $dm$  is a tuple  $lcm = (olc, olc_{DEP})$  where*

- $olc : OT \mapsto \mathcal{OLC}$  assigns to each object type  $ot \in OT$  (of the data model) an object life cycle  $olc(ot) \in \mathcal{OLC}$
- $olc_{DEP} : RT \mapsto \mathcal{OLC}_{DEP}$  assigns to each relation type  $rt = (ot1, rt, ot2) \in RT$  an OLC dependency  $olc_{DEP}(rt)$  for the object life cycles  $olc(ot1)$  and  $olc(ot2)$  of the object types  $ot1$  and  $ot2$ .

Regarding the execution of created process structures, it is important to guarantee soundness, i.e., to ensure that data-driven process structures terminate with a correct end state. Deadlocks might occur (1) when external state transitions are starting from non-deterministic states, and (2) when external state transitions are forming cycles. The first situation can be avoided during runtime, for example, using deadpath elimination techniques. The second situation can be recognized during buildtime by analyzing the process structure. Analyzing large data-driven process structures, however, generates high efforts. COREPRO enables checking soundness on model level and guarantees soundness for every data-driven process structure that bases on a sound LCM.

To check soundness of a LCM, all OLCs and OLC dependencies (i.e., their external state transitions) of the LCM are composed. In addition, a unique start state is added and connected with all start states of the OLCs via *direct state transitions* (cf. Fig. 6). Accordingly, all OLC end states are connected with a unique end state. Direct state transitions are concurrently triggered and lead to the deactivation of the source state and the activation of the target state (cf. Table 1). We denote the extended transition system resulting from this as *LCM machine*. The LCM machine constitutes a LCS for a data structure where every element and relation type is instantiated once (to map OLC dependencies for recursive relation types, it is necessary to contemplate two OLCs for the object type associated with the recursive relation type). Thereby, efforts for soundness checks do not rise with the size of the instantiated process structure but only depend on the size of the LCM machine. As example consider Fig. 6, which shows the LCM machine for the LCM depicted in Fig. 5a. Since sub-processes connected with state transitions do not affect soundness checks (we presume sound sub-processes), they can be neglected when analyzing soundness.

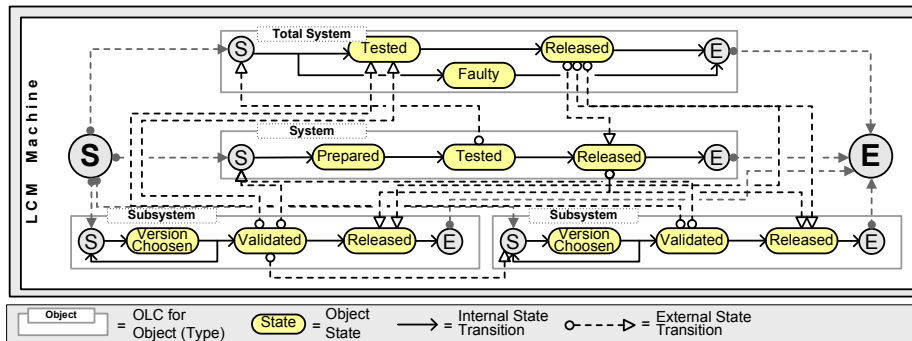


Fig. 6. LCM Machine for the LCM from Fig. 5a

**Definition 6 (Soundness of the LCM).** A LCM machine is sound if each state of the LCM machine (including its end state) can be enabled by direct or internal state transitions beginning with the start state of the LCM machine and every external transition becomes activated (or deactivated) then.

### 4.3 Creating the Life Cycle Coordination Structure

So far, we have introduced the data part which comprises the data model and the data structure, and the LCM (consisting of OLCs and OLC dependencies) which integrates the data model and the sub-processes. Based on this, the data-driven process structure, i.e., the *life cycle coordination structure* (LCS), can be automatically derived for respective data structures. The LCS comprises a start and an end state, an OLC instance for every object in the data structure, and external state transitions between these OLC instances according to the relations defined between the objects (cf. Fig. 5b).

**Definition 7 (Life Cycle Coordination Structure).** Let  $dm = (O, R)$  be a data structure and let  $lcm = (olc, olc_{DEP})$  be a life cycle coordination model. Then: A life cycle coordination structure based on  $dm$  and  $lcm$  is a tuple  $lcs = (olc_{inst}, est_{inst}, s_{start}, s_{end}, ST, ET)$  where

- $olc_{inst} : O \mapsto \mathcal{OLC}$  assigns to each object  $o \in O$  an instance of the associated object life cycle  $olc_{inst}(o) = olc(objtype(o))$
- $est_{inst} : R \mapsto \mathcal{OLC}_{DEP}$  assigns to each relation  $r \in R$  the associated external state transitions  $est_{inst}(r) = est(olc_{DEP}(reltype(r)))$
- $s_{start}$  denotes the initial state and  $s_{end}$  the final state
- $ST$  is the set of direct state transitions connecting the start state of the LCS with the start states of the instantiated object life cycles
- $ET$  is the set of direct state transitions connecting the end states of the instantiated object life cycles with the end state of the LCS.

The operations for creating a LCS are defined in Table 2. Based on a data structure and an LCM, three steps become necessary to generate the LCS. Algorithm 1 describes these steps in detail:

1. For every object in the data structure, the OLC associated with the corresponding object type is instantiated.
2. For every relation in the data structure, the OLC dependencies associated with the corresponding relation type (i.e., their external state transitions) are inserted to connect states of the dependent OLCs.
3. Direct state transitions are inserted, which connect the LCS start state with all OLC start states and all OLC end states with the LCS end state.

As result, we obtain the complete LCS representing the logical view on the data-driven process structure (cf. Fig. 5b). Such LCS can be transformed to activity-centered process representations, like BPMN or WS-BPEL.

Checking soundness of an LCS comprising hundreds up to thousands of sub-processes and (external) state transitions is a complex task to accomplish. COREPRO enables soundness checking on model level and ensures that every LCS created on basis of a sound LCM is sound as well (cf. Theorem 1).

**Theorem 1 (Soundness of the LCS).** *Assume that an LCS has been created with Alg. 1 with a particular data structure and an LCM as input. Then: If the LCM machine of the LCM is sound (cf. Definition 6), the created LCS is sound as well.*

**Table 2.** Operations for Creating an LCS

Operation	Effect
createLCS	Creates a new LCS
insertStartState(lcs)	Inserts the initial state $s_{start}$ to the given lcs
insertEndState(lcs)	Inserts the final state $s_{end}$ to the given lcs
insertOLC(lcs,olc)	Inserts an instance of the given OLC to the given lcs
insertExtTrans(lcs,s,p,s')	Inserts an external state transition from state $s = (\text{Transitionsystem}, \text{State})$ to state $s' = (\text{Transitionsystem}, \text{State})$ with the associated sub-process $p$ to the given lcs
insertDirTrans(lcs,s,s')	Inserts a direct state transition from state $s = (\text{Transitionsystem}, \text{State})$ to state $s' = (\text{Transitionsystem}, \text{State})$ to the given lcs

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1 Input:  $DS = (O, R)$ ,  $LCM = (OLC, OLC_{DEP})$ 
2 Output:  $lcs = (OLC_{inst}, EST_{inst}, s_{start}, s_{end}, ST, ET)$ 
3 // Initialize the LCS and insert start and end state
4  $lcs := \text{createLCS}$ ;  $s := \text{insertStartState}(lcs)$ ;  $e := \text{insertEndState}(lcs)$ ;
5 // Insert an OLC instance for every instantiated object and connect it with the start
  and end state of the LCS via directed state transitions
6 forall  $obj \in O$  do
7    $olc := \text{insertOLC}(lcs, olc(objtype(obj)))$ ;
8    $\text{insertDirTrans}(lcs, (lcs, s_{start}), (olc, s_{start}(olc)))$ ;
9    $\text{insertDirTrans}(lcs, (olc, s_{end}(olc)), (lcs, s_{end}))$ ;
10 // Insert external state transitions for each instantiated relation
11 forall  $rel = (o_1, rt, o_2) \in R$  do
12   // Insert external state transitions between the OLC of  $o_1$  and the OLC of  $o_2$ 
13   forall  $est = (s_1, p, s_2) \in olc_{DEP}(rt)$  do
14      $\text{insertExtTrans}(lcs, (olc_{inst}(o_1), s_1), p, (olc_{inst}(o_2), s_2))$ ;
15 return(lcs);

```

**Algorithm 1:** Generation of the Life Cycle Coordination Structure

Theorem 1 can be inductively proven: The LCM machine constitutes a sound LCS for a data structure where every object type and every relation type are instantiated once. When adding one additional object and all corresponding relations to other objects, we can prove that this leads to a sound LCS again ( $n = 1$ ). Making this assumption for  $n$  additional objects, we can show that soundness can still be guaranteed when adding a further object and corresponding relations ( $n \rightarrow n + 1$ ). Due to lack of space, we omit further details.

#### 4.4 Change Scenarios

When dealing with data-driven process structures, change management becomes an important issue. Adaptations of process structures become necessary, for example, when the underlying data structure is changed (e.g., when adding a new object). In COREPRO, such changes can be specified at the data level and are then automatically translated into corresponding adaptations of the process structure. Compared to conventional approaches (e.g., the manual adaptation of activity-centered process structure representations), adaptation efforts can be significantly reduced. To illustrate this, we sketch three change scenarios in which users (e.g., engineers) adapt the (product) data structure.

##### *Removing an object*

*Example:* The **Speed Sensor** subsystem shall not be processed any longer, i.e., the **Speed Sensor** and all its relations to or from other objects are removed from the data structure (cf. Fig. 3b).

*Conventional approach:* Manually removing associated sub-processes and their incoming and outgoing synchronization links from the process structure.

*COREPRO procedure:* The corresponding OLC and the external state transitions are automatically removed from the LCS.

##### *Removing a relation*

*Example:* The **Main Unit** does not use the **Speed Sensor** any longer, i.e., the relation **usesSubsystem** between the **Main Unit** and the **Speed Sensor** object is removed from the data structure (cf. Fig. 3b).

*Conventional approach:* Manually removing the sub-process dependencies which are no longer necessary from the process structure (i.e., synchronization links of the sub-processes modifying the **Main Unit** and the **Speed Sensor** subsystem).

*COREPRO procedure:* The corresponding external state transitions are automatically removed from the LCS.

##### *Adding an object*

*Example:* A new **Head-Up Display** subsystem shall be processed as part of the **Navigation** system, i.e., a new object is added to the data structure and related to the **Navigation** object (cf. Fig. 3b).

*Conventional approach:* Manually inserting the corresponding sub-processes and associated synchronization links.

*COREPRO procedure:* The corresponding OLC and the external state transitions are automatically added to the LCS.

Taking these scenarios, COREPRO allows users to apply process structure changes at a high level of abstraction. Using conventional, activity-driven approaches for process modeling, process structures would have to be manually adapted in case of a data structure change. That requires extensive process knowledge and necessitates additional soundness checks. By contrast, COREPRO significantly reduces efforts for adaptation. The process structure can be adapted without comprehensive process knowledge by simply changing the data structure. The soundness of the resulting process structure is assured, since the model level (i.e., the data model and the LCM) remains unchanged.

## 5 Practical Impact

To indicate the practical benefit of our modeling approach, we introduce a model calculation for a characteristic process from the car development domain: the release management (RLM) process for electrical systems [1, 3]. The calculation (cf. Table 3) bases on the experiences we gained from case studies in this domain.

The release of an electrical system encompasses 200 to 300 components (depending on the car series), which are divided in root components and their variants (e.g., driver's airbag as root component and passenger's airbag as its variant). Root components are further divided into categories requiring different processing. For example, releasing a multimedia component requires different sub-processes when compared to a security related component. Additionally, components are grouped into systems (e.g., navigation system covers several components) to integrate logically coherent components. Finally, systems are collected in total systems, which represent the car series to be developed (e.g., S-Class). Altogether, this leads to the definition of a data model with about 20 object types and 25 relation types connecting them.

On instance level, the relation types with a 1:n cardinality lead to more than 200 instantiated relations. Additionally, there exist dependencies between components that exchange signals and messages. These relation types are defined with an n:m cardinality leading to more than 400 relations. OLCs for the different object types coordinate 5 (for components) to 20 (for systems) sub-processes.

In contrast to conventional modeling, the creation of the process structure constitutes an automated task. In total, the LCS contains 200 to 300 OLCs (according to the number of objects) with more than 1300 sub-processes. Relation types encompass 2 to 6 external state transitions (cf. Table 3) leading to more than 1500 external state transitions within the generated LCS.

The potential for reducing modeling efforts when using the instantiation mechanism of COREPRO depends on the ratio of object types to objects. Even though the calculation bases on a moderate estimate, it indicates that the modeling efforts for RLM processes can be significantly reduced by more than 90%. The benefit even increases considering the fact that soundness checks (cf. Sect. 4.2) and changes (cf. Sect. 4.4) become less complex.

**Table 3.** Projection of the Modeling Efforts Reduction

Model			
Data Model		LCM	
Object Types	Relation Types	Sub-Processes	Ext. Transitions per Relation Type
20	25	5-20	2-6
Instance			
Data Structure		LCS	
Objects	Relations	Sub-Processes	Ext. Transitions
200-300	>600	>1300	>1500
Modeling Efforts Reduced by			
>90%	>95%	>98%	>99%

## 6 Related Work

Activity-driven approaches for process modeling do not focus on the automated creation of large process structures comprising multiple sub-processes. Interaction graphs [15] and choreography definition languages (e.g., [16]), for example, are activity-centered approaches to specify the choreography between distributed sub-processes. The data-driven derivation of sub-process dependencies is not covered. Other approaches partially enable the data-driven creation of process structures. For example, multiple instantiation of activities based on simple data structures (set, list) [17, 18] is supported by UML 2.0 activity diagrams (*Expansion Region*) [14] and BPMN (*Multiple Instances*) [19]. They enable iterated or concurrent execution of the same activity for each element given by a flat data container. Utilization of respective data structures raises further options, such as data-driven process control with exception handling mechanisms [20]. However, these approaches aim at the sequential or concurrent execution of multiply instantiated activities. COREPRO, by contrast, focuses on the definition of arbitrary complex data structures and their association with (sub-)processes. The data-driven process structure in Fig. 7, for example, realizes the interleaved synchronization of sub-processes. It represents a list structure where sub-processes are executed for every list element. In this context, COREPRO also allows for the realization of anticipation concepts [21]: Regarding the generated LCS, **Process A** for **Element 2** can be started even though the processing of **Element 1** has not been finished.

Approaches for explicitly generating process structures based on bills of material are described in [5, 7]. The idea of coordinating activities based on data dependencies also constitutes the basis of the *Case Handling* paradigm [6]. The idea is to model a process structure by relating activities to the data flow. The concrete activity execution order at runtime then depends on the availability of data. Another approach integrating control and data flow is provided by *AHEAD* [22], which offers dynamic support for (software) development process structures. The approach enables the integration of control and data flow, by relating activities to the objects defined in the data model. Based on this information, dynamic task nets are generated. The goal of respective data-driven approaches is the precise mapping of object relations to sub-process dependencies. Though the object relations indicate sub-process dependencies (cf. Sect. 2), the informa-



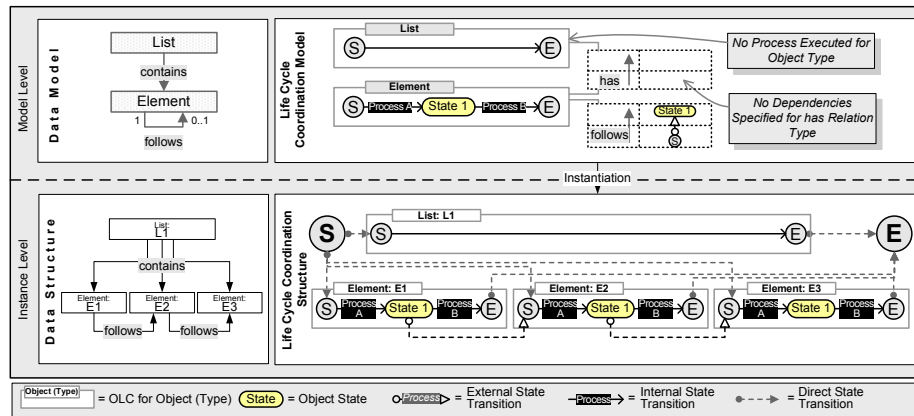


Fig. 7. Interleaved Processing of a List

tion given by relations is insufficient for their direct mapping to synchronization edges for three reasons. First, *several* sub-processes may modify one object, whereas the relation itself does not reflect which of these sub-processes have to be synchronized. Second, the relations do not provide sufficient information about the direction of synchronization dependencies. In Fig. 1, for example, the relation points from the *Speed Sensor* to the *Navigation* object, while the synchronization dependencies between their sub-processes point in both directions. Third, it is also requested to associate synchronization dependencies with the enactment of synchronizing sub-processes (cf. Sect. 4.2).

A general approach following the idea of modeling life cycles and relating them is *Object/Behavior Diagrams*. The concept allows for the object-oriented definition of data models which can be enhanced by runtime aspects [8]. The behavior is defined for every object within a Petri Net relied life cycle diagram. Another approach using life cycles for describing operational semantics of business artifacts (semantic objects) is *Operational Specification* (OpS) [11]. The collection of all objects and their life cycles specifies the operational model for the entire business. The *Object-Process Methodology* (OPM) is an object-oriented approach from the engineering domain. It focuses on connecting objects (or object states) and processes by procedural links [10]. *Team Automata* provide a formal method to describe the connection of labeled transition systems (automata) via external actions associated with (internal) transitions [9]. Automata including transitions with the same external action perform them simultaneously. The idea is adopted in [23] where Team-Automata are structured in an object-oriented way. Its synchronization mechanisms are based on events. These approaches rather focus on an activity-driven specification of dependencies (based on events) than on the consideration of data relations for process structure generation. In contrast to event-based synchronization, external state transitions can be added, removed or disabled (e.g., in order to avoid deadlocks) without changing the dynamic aspects of the object itself (i.e., the OLC).

## 7 Summary and Outlook

The COREPRO approach offers promising perspectives with respect to the modeling and coordination of large data-driven process structures consisting of numerous sub-processes and their interdependencies. COREPRO supports the loose coupling of data and sub-processes by defining life cycles of data objects, and it provides a well-defined mapping of object relations to OLC dependencies. Further, COREPRO distinguishes between model and instance level, which enables a high level of abstraction, extensibility and reuse. In particular, modeling and change effectiveness are significantly enhanced by

- introducing model-driven design in conjunction with an instantiation mechanism for data-driven process structures
- enabling the instantiation of different data structures and automatically generating respective data-driven process structures
- integrating data and processes which allows users without process knowledge to adapt the process structures by changing the data structure.

Another important issue to be considered is the need for flexibility at runtime, such as applying structural changes during enactment (cf. Sect. 4.4). This becomes necessary, for example, when the number of objects or relations between them is not (exactly) known at buildtime. Due to the many sub-process dependencies, uncontrolled runtime changes may lead to inconsistencies not only within single OLCs, but also within the whole LCS. In addition to structural changes, we also have to consider state changes. To realize iterative development processes, for example, data structures need to be processed several times. However, that necessitates the (partial) utilization of previous processing states of objects; i.e., object states which were already activated before execution, have to be retained. For example, product components which have already been tested and which remain unchanged do not need to be tested again. Applying such changes and supporting exceptional situations (e.g., abnormal termination of a sub-process or backward jumps within an OLC) while preserving consistency is a challenging problem [2]. Solutions for runtime scenarios and exception handling are also addressed by COREPRO and will be presented in future publications.

We have implemented major parts of the presented modeling concepts in a prototype, which we use for a first proof-of-concept case study in car development [24]. The approach will be applied for modeling, coordinating and maintaining data-driven process structures in the automotive industry. However, the presented concept is not specific to the engineering domain. We also plan to evaluate COREPRO in the healthcare domain, where the approach shall be used to model medical treatment processes [25].

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